

RUSTHORNBELT:

A Semantic Foundation for Functional
Verification of **Rust** Programs with Unsafe Code



Yusuke Matsushita

The University of Tokyo

Jacques-Henri Jourdan

CNRS, LMF

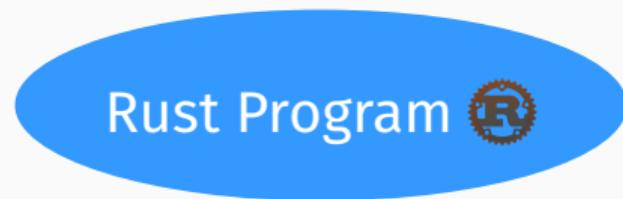
Xavier Denis

Université Paris-Saclay, LMF

Derek Dreyer

MPI-SWS

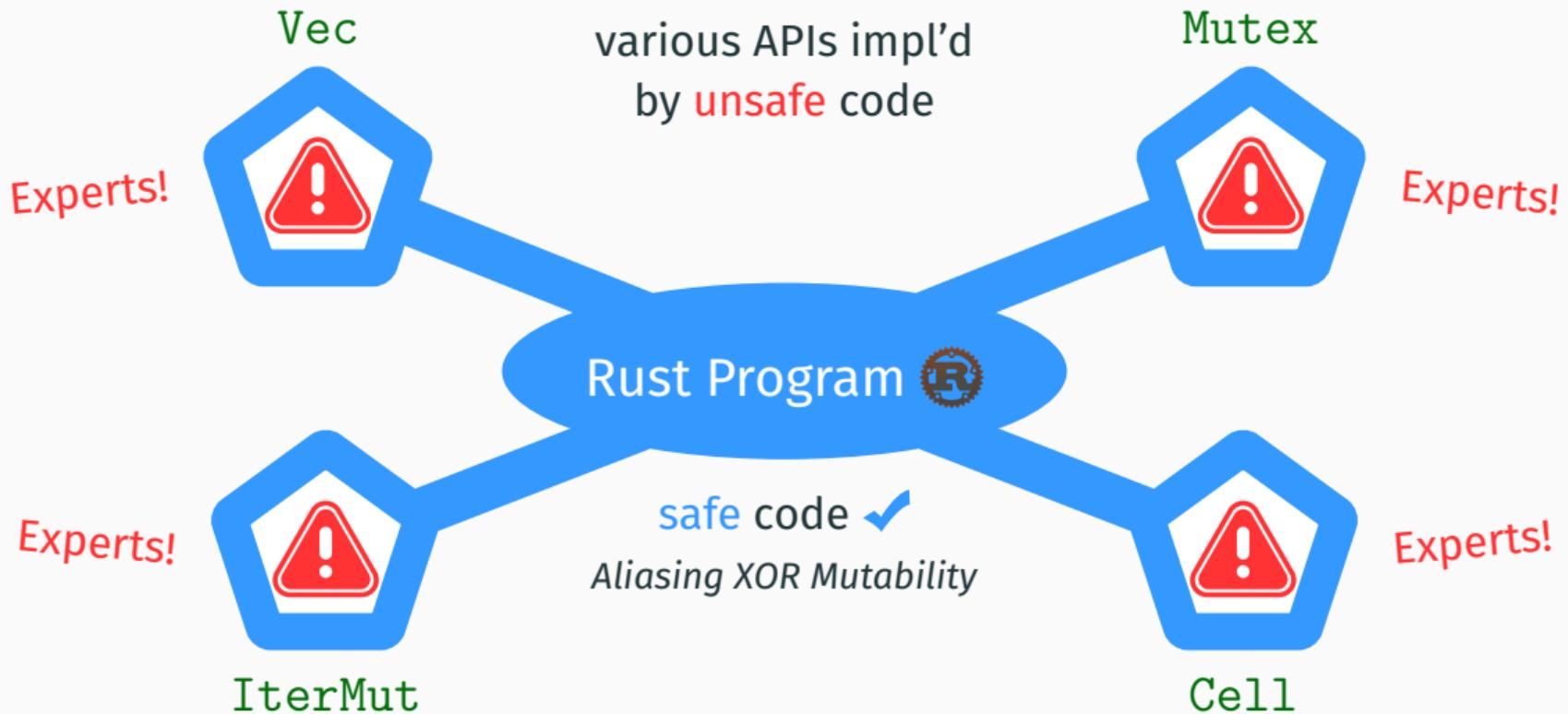
Our High-level Goal: Rust Verification



safe code ✓

Aliasing XOR Mutability

Our High-level Goal: Rust Verification



Two Lines of Prior Work on Rust Verification



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RUSTBELT 

[Jung+ '18]

Safe & Unsafe 

Manual in IRIS *

Type Safety

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RUSTHORN 

[Matsushita+ '20],

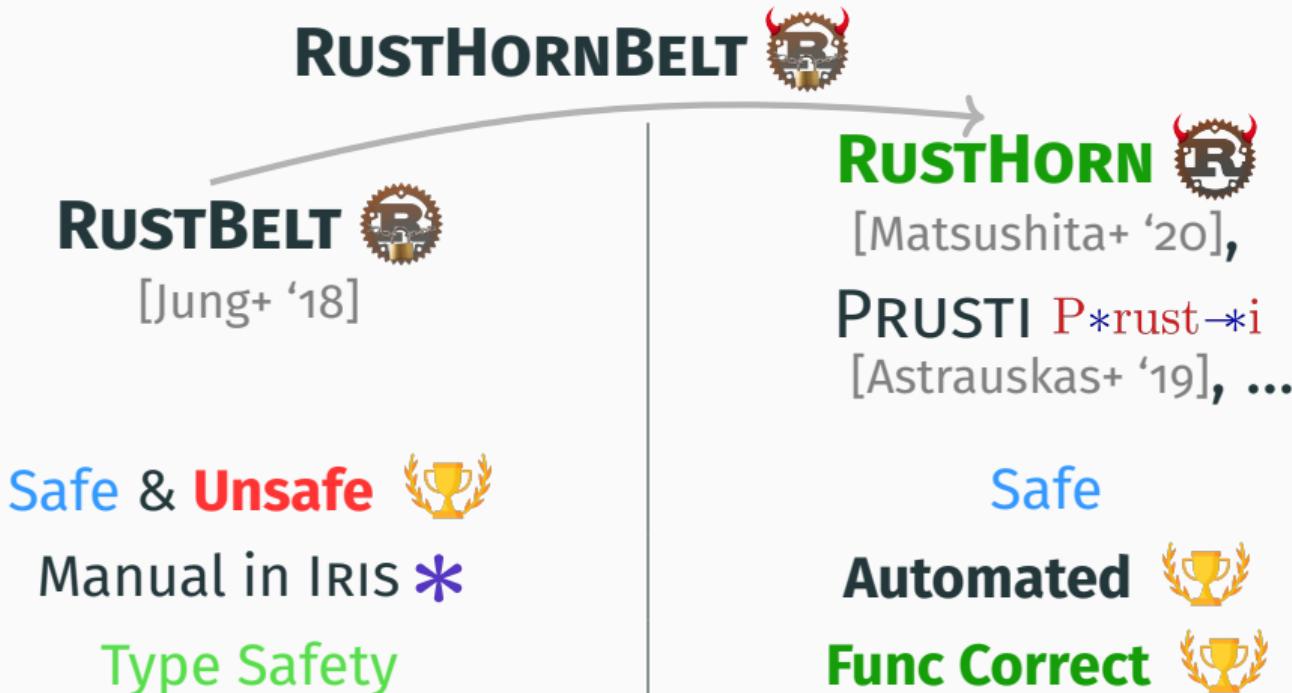
PRUSTI P_{rust}-*i
[Astrauskas+ '19], ...

Safe

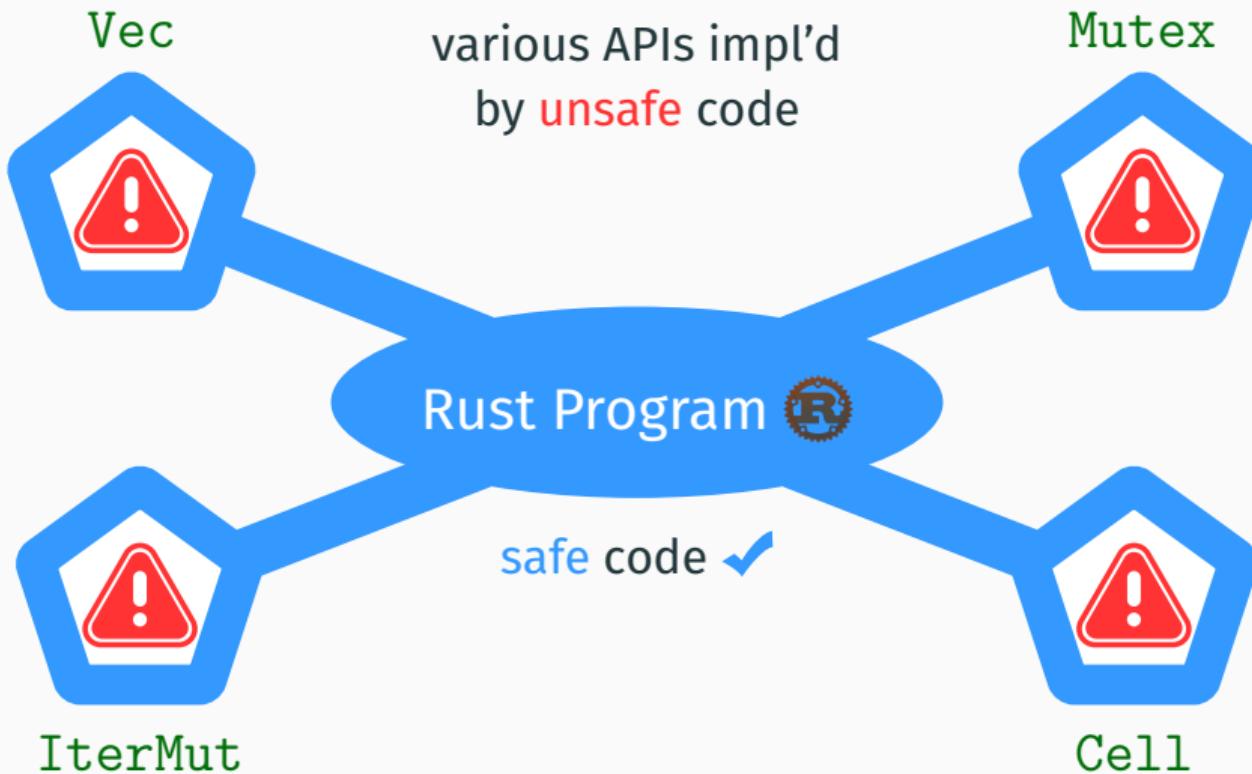
Automated 

Func Correct 

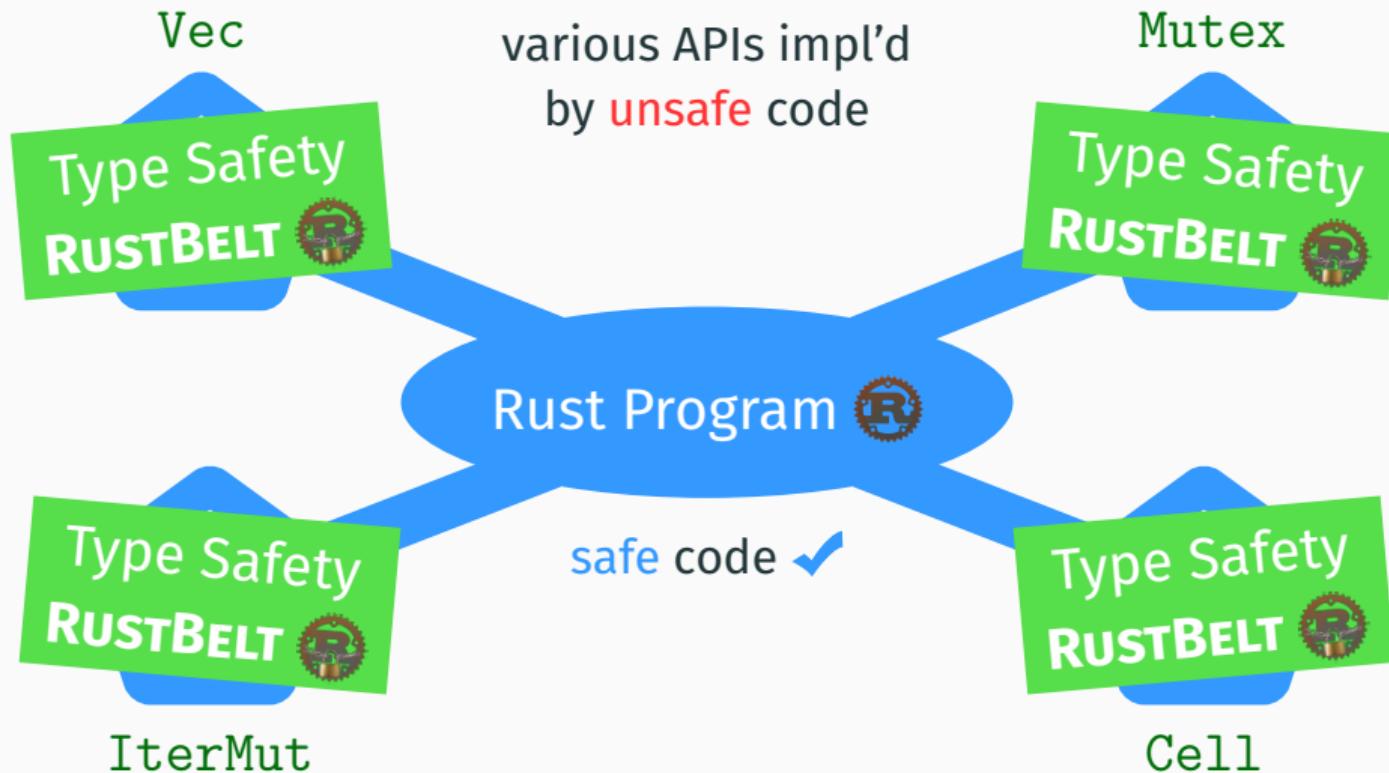
Two Lines of Prior Work on Rust Verification



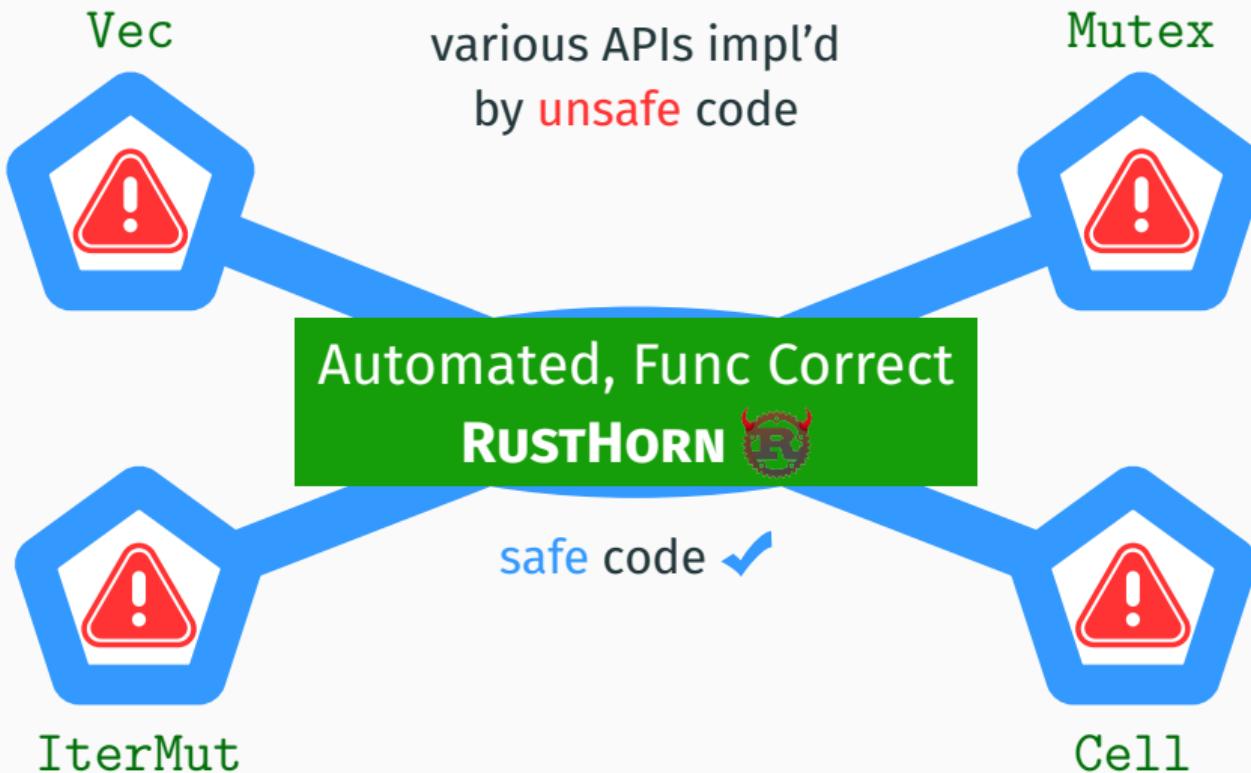
Our Goal: Marrying RUSTBELT with RUSTHORN



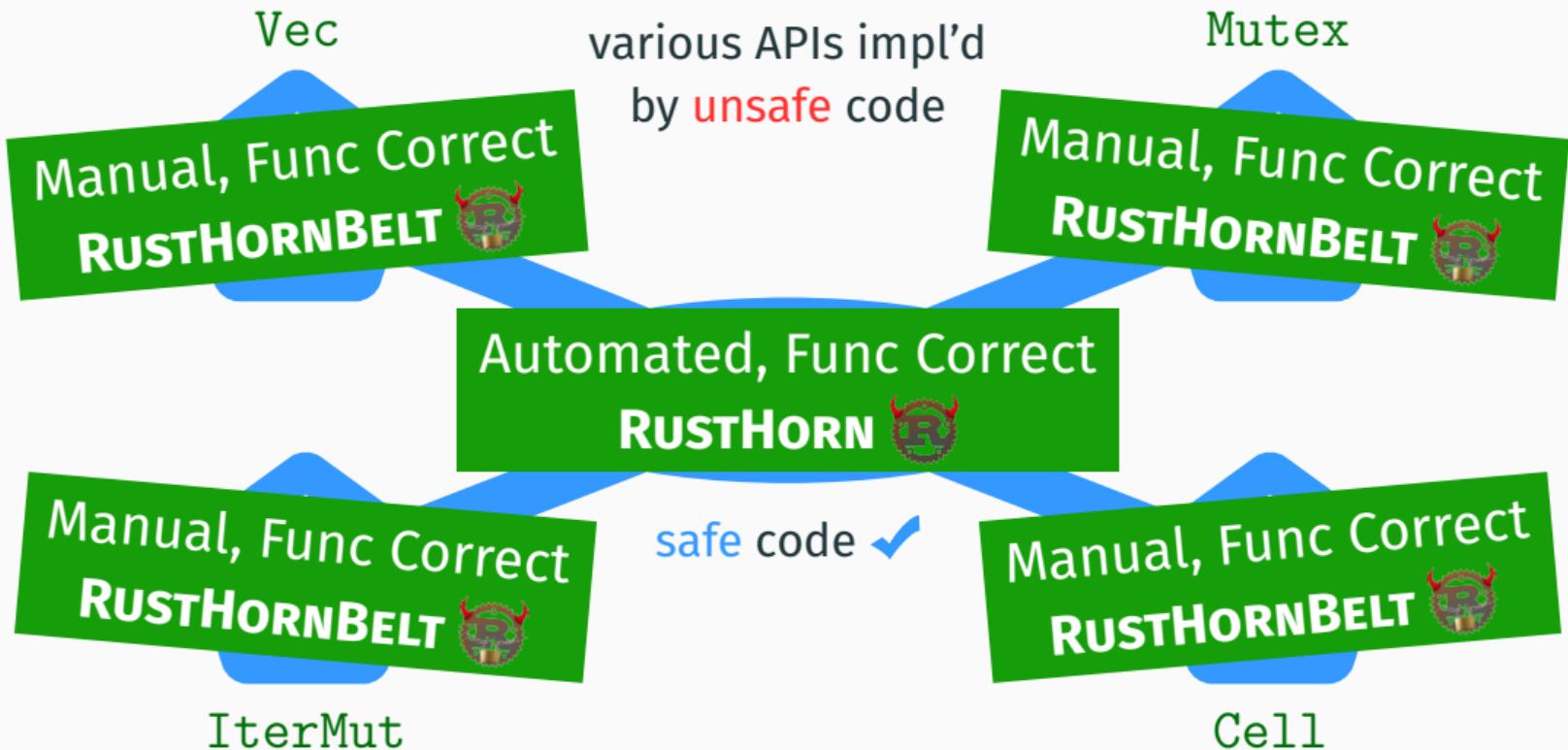
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Our Work, RUSTHORNBELT



Marriage of **RUSTBELT**  with **RUSTHORN** :

Our Work, RUSTHORNBELT



Marriage of **RUSTBELT**  with **RUSTHORN** :

- **RUSTBELT**-based **semantic model** of Rust's **ownership types** in the separation logic **IRIS *** (Coq !), but extended with **RUSTHORN**-style functional **specs**

Our Work, RUSTHORNBELT



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Our Work, RUSTHORNBELT



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Motivating Example for RUSTHORN

What spec do you give to `max_or_incr`?

```
fn max_or_incr(ma: &mut int, mb: &mut int) -> &mut int {  
    if *ma >= *mb {  
        *mb += 42;    ma  
    } else {  
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fn max_or_incr(ma: &mut int, mb: &mut int) -> &mut int {  
    if *ma >= *mb {  
        *mb += 42 · ma  
    } else {  
        *ma -> Mutable reference,  
borrowing ownership  
    }  
}
```

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fn max_or_incr(ma: &mut int, mb: &mut int) -> &mut int {  
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        *mb +: Not aliased }  
    } else {  
        *ma += 42; mb  
    }  
}
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Separation Logic?

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~~Separation Logic?~~

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- Key idea: **Mutable ref** `ma` repr'ed as a **value** (a , a')
 - a — *current value*
 - a' — **prophecy**  of *final value*, returned to original owner

RUSTHORN's Answer to the Example

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        *mb += 42;    ma  
    } else {  
        *ma += 42;    mb  
    }  
}
```

$$\begin{aligned} \text{max_or_incr } (a, a') (b, b') \text{ res} &\triangleq \\ \text{if } a \geq b \text{ then } b' = b + 42 \wedge \text{res} = (a, a') \\ \text{else } a' = a + 42 \wedge \text{res} = (b, b') \end{aligned}$$

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 b then $\underline{b'} = b + 42 \wedge res = (a, \underline{a'})$
else $\underline{a'} = a + 42 \wedge res = (b, \underline{b'})$

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Type-Spec System

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RustBelt's type system

$$T \vdash I \dashv r. T'$$

typing judgment

Type-Spec System

RustBelt's type system + Func Spec

Type-Spec System

RustBelt's type system + Func Spec

predicate transformer

$|\mathbf{T}'| \rightarrow \text{Prop} \ni \text{post} \mapsto \text{pre} \in |\mathbf{T}| \rightarrow \text{Prop}$

RUSTHORN-style Representation of Rust Objects

Each **object** of Rust type $\textcolor{green}{T}$ is represented by RUSTHORN
as a **value** of *logic sort* $[\textcolor{green}{T}]$

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$$[\textcolor{blue}{\text{int}}] \triangleq \mathbb{Z} \quad [\textcolor{green}{(T, U)}] \triangleq [\textcolor{green}{T}] \times [\textcolor{green}{U}]$$

$$[\&\textcolor{blue}{\text{mut }} T] \triangleq [T] \times \underline{[T]}$$

prophecy


Interpreting Rust's Types & RUSTHORN in IRIS

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Semantics of the typing judgment in IRIS \ast , roughly:

$$\llbracket T \vdash I \dashv r. T' \rrbracket \triangleq \{ \llbracket T \rrbracket \} / \{ r. \llbracket T' \rrbracket \}$$



Interpreting Rust's Types & RUSTHORN in IRIS

Semantics of the type-spec judgment in IRIS \ast , roughly:

$$[\![\mathbf{T} \vdash I \dashv r. \mathbf{T}' \rightsquigarrow \Phi]\!] \triangleq \forall post.$$

$$\{ \exists \bar{a}. \Phi post \bar{a} * [\![\mathbf{T}]\!](\bar{a}) \} \mid \{ r. \exists \bar{b}. post \bar{b} * [\![\mathbf{T}']\!](\bar{b}) \}$$



Example — Rust's Vec API

`Vec<T>` — Growable, heap-allocated array

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fn push(mv: &mut Vec<T>, a: T)
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```
fn index_mut(mv: &mut Vec<T>, i: int) -> &mut T
```

$\lfloor \text{Vec} < \text{T} \rfloor \triangleq \text{List } \lfloor \text{T} \rfloor$

RUSTHORN's Spec for push

`mv: &mut Vec<T>, a: T ⊢ push(mv, a) ⊢`

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RUSTHORN's Spec for index_mut

```
mv: &mut Vec<T>, i: int ⊢ index_mut(mv, i) ⊢ ma. ma: &mut T
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`mv: &mut Vec<T>, i: int ⊢ index_mut(mv, i) ⊢ ma. ma: &mut T`

$\rightsquigarrow \lambda post, ((v, v'), i). \quad 0 \leq i < |v| \wedge$

...

RUSTHORN's Spec for index_mut

$\text{mv: } \&\text{mut Vec<T>, i: int} \vdash \text{index_mut}(\text{mv}, \text{i}) \dashv \text{ma. ma: } \&\text{mut T}$

$\rightsquigarrow \lambda post, ((v, v'), i). \quad 0 \leq i < |v| \wedge$

$v' = v\{i := ?\} \rightarrow post(v[i], ?)$

resolution

RUSTHORN's Spec for index_mut

`mv: &mut Vec<T>, i: int ⊢ index_mut(mv, i) ⊢ ma. ma: &mut T`

$\rightsquigarrow \lambda post, ((v, v'), i). \quad 0 \leq i < |v| \wedge$

$\forall a'. \underline{v' = v\{i := a'\}} \rightarrow post(v[i], \underline{a'})$

resolution

prophecy



Parametric Prophecies

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Prior work [Jung+ '20] — *Not flexible enough*

Parametric Prophecies

RustHorn  – **Mutable ref** as (a, a') , a' is **prophecy** 

← Semantically model this in separation logic IRIS?

Prior work [Jung+ '20] – *Not flexible enough*

Our solution  – **Clairvoyant monad** $\text{Clair } A \triangleq \text{Future} \rightarrow A$

i.e., reader monad over every possible **future** π

Semantics in IRIS – Now Prophetic!

$$\llbracket \mathbf{T} \vdash I \dashv \mathbf{r}. \mathbf{T}' \rightsquigarrow \Phi \rrbracket \triangleq \forall post.$$

$$\{ \exists \bar{a}. \quad \Phi \text{ } post \quad \bar{a} \quad * \llbracket \mathbf{T} \rrbracket(\bar{a}) \}$$

$$I \quad \{ \mathbf{r}. \exists \bar{b}. \quad post \quad \bar{b} \quad * \llbracket \mathbf{T}' \rrbracket(\bar{b}) \}$$

Semantics in IRIS – Now Prophetic!

$$\llbracket \mathbf{T} \vdash I \dashv \mathbf{r}. \mathbf{T}' \rightsquigarrow \Phi \rrbracket \triangleq \forall \hat{post}.$$

$$\{ \exists \bar{\hat{a}}. \langle \lambda \pi. \Phi (\hat{post} \pi) \bar{\hat{a}} \pi \rangle * \llbracket \mathbf{T} \rrbracket (\bar{\hat{a}} \pi) \}$$

$$I \quad \{ \mathbf{r}. \exists \bar{\hat{b}}. \langle \lambda \pi. \hat{post} \pi (\bar{\hat{b}} \pi) \rangle * \llbracket \mathbf{T}' \rrbracket (\bar{\hat{b}} \pi) \}$$

In the Paper, Also...

- More on type-spec system &
key Rust APIs (`Vec`, `IterMut`, `Mutex`, ...) (§2)
- More on **parametric prophecies**  &
Model of **mutable ref** `&mut T` &
Proof sketch of key type-spec rules (§3)
- **Coq**  **mechanization** details &
Automation **benchmarks** in CREUSOT (§4)

